



The Trend Locality Sensitive Hash: TLSH

Jon Oliver, Chun Cheng, Yanggui Chen

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Contact: tlsh@trendmicro.com

Getting TLSH

- Contact us at
tlsh@trendmicro.com
- Source Code:
<https://github.com/trendmicro/tlsh/>

What is Locality Sensitive Hashing

- Traditional hashes (such as SHA1 and MD5) have the property that a small change to the file being hashed results in a completely different hash
- Locality Sensitive Hashes (LSH) have the property that a small change to the file being hashed results in a small change to the hash
 - You can measure the similarity between 2 files by comparing their LSH values

Example Locality Sensitive Hashing



Text 1 – Chapter 1 of Pride and Prejudice

It is a truth universally acknowledged, that a single man in possession of a good fortune, must be in want of a wife.

..

When she was discontented, she fancied herself nervous.
The business of her life was to get her daughters married; its solace was visiting and news.

Text 2 - Chapter 1 of Pride and Prejudice with last line removed

It is a truth universally acknowledged, that a single man in possession of a good fortune, must be in want of a wife.

..

When she was discontented, she fancied herself nervous.

Example Locality Sensitive Hashing

TLSH

Text 1 E491A51FA380022245B021E9770F3A6FF706C1780365C631581EF6263731EAA87F96EE

Text 2 5B91940FA380026245B021A9771F7A6FF706C1780765C671981EF6263731EAA87F96DE

MD5_HASH

Text 1 3b9dd1f86ce0c3b467055b48f9a5221c

Text 2 7dc8267c6bea14d36df64934aad4604f

SHA1_HASH

Text 1 8b8c6ce1253515a1fbceaec0f5cfc58780e6fd5e

Text 2 e494d7fa7b4080520c59a6702764983ff9b6d399

The MD5 and SHA1 hashes are completely different

For these 2 pieces of text, the TLSH values are quite similar Hashes

Example Locality Sensitive Hashing

TLSH

Text 1 E491A51FA380022245B021E9770F3A6FF706C1780365C631581EF6263731EAA87F96EE

Text 2 5B91940FA380026245B021A9771F7A6FF706C1780765C671981EF6263731EAA87F96DE

The distance between Text 1 and Text 2

$$\text{distance}(\text{Text1}, \text{Text2}) = 11$$

Distance scores can go up to 1000 and above

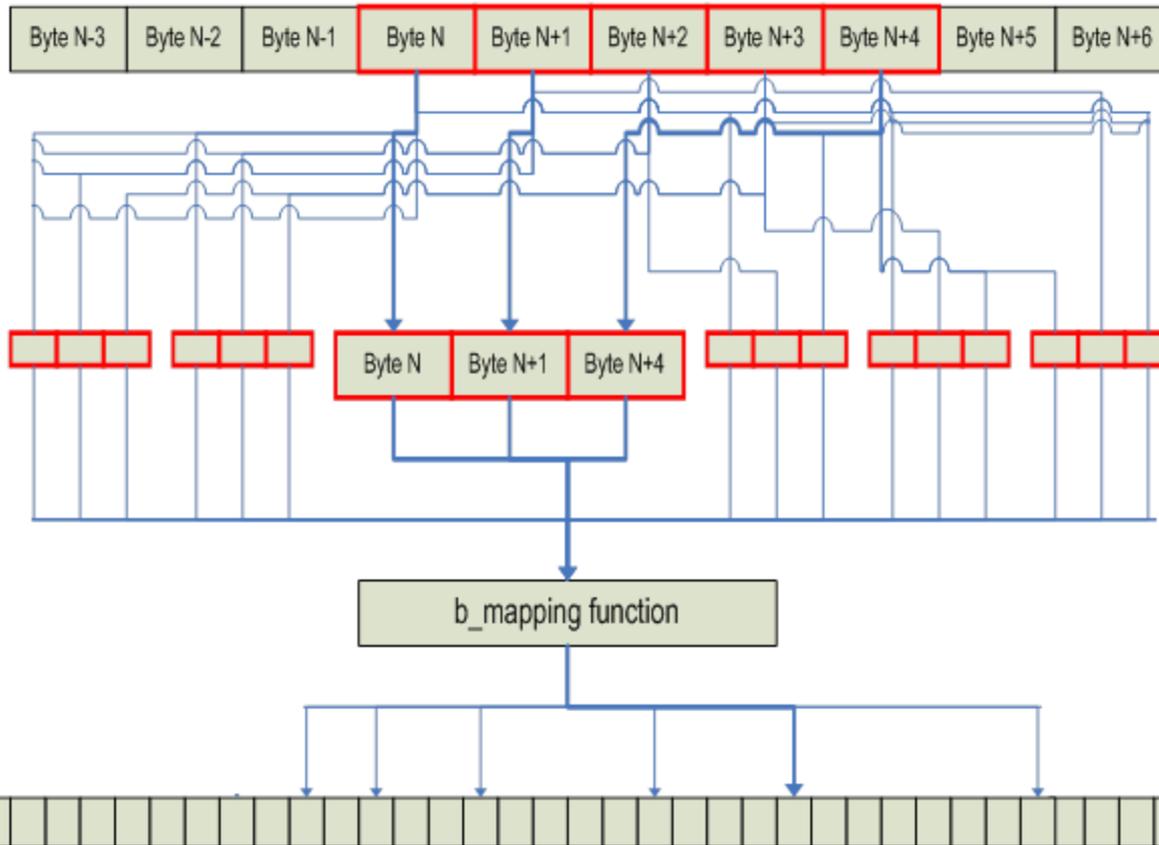
A low score (of 50 or less) means that the files are quite similar

You will need to determine an appropriate threshold for your application

A distance of 0 means that the files are (very likely) to be exactly the same

Just like the MD5 and SHA1 schemes, collisions can occur and very different files will have the same hash value.

Algorithm to determine TLSH



Each sliding window contains 5 bytes

Those 5 bytes form trigrams

b_mapping function maps each trigram into one of the buckets

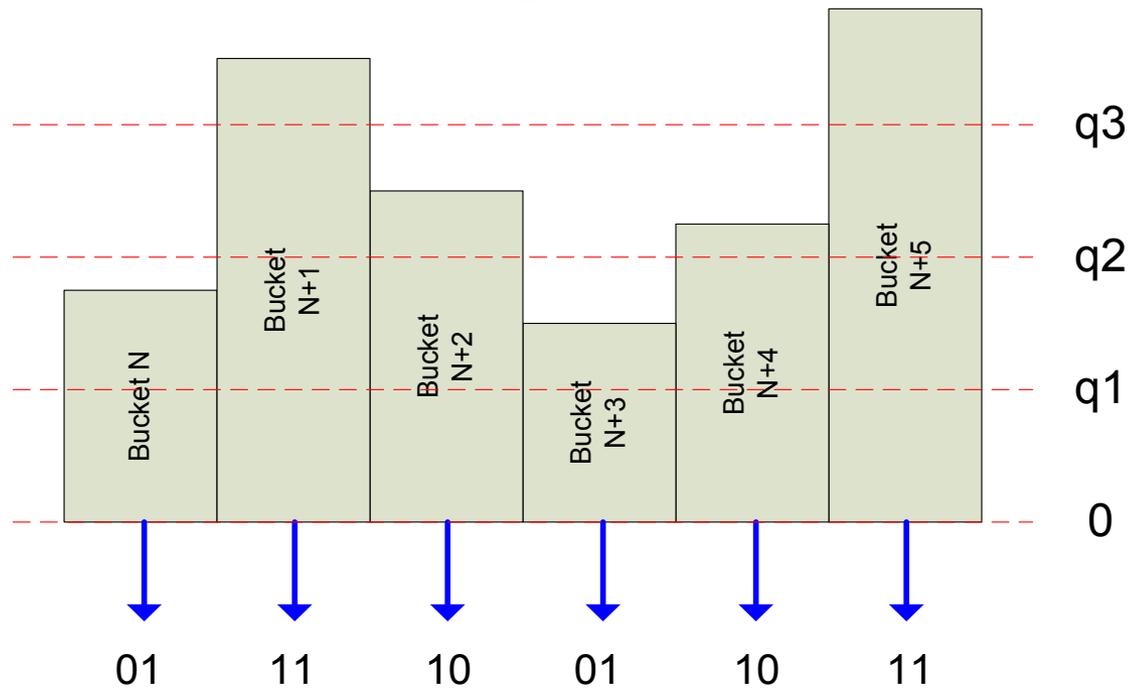
For each mapping, increase the bucket count by 1

Algorithm to determine TLSH

- We use the Pearson hash [reference 1] as the mapping function between the trigrams from a window to the buckets.

Algorithm to determine TLSH

- TLSH uses a 4-way to reflect the differences between different histograms
- The q2 point is at the median bucket count
- The q1 are the lower and higher quartiles respectively



Algorithm to determine the hash

- Introduce three head bytes (6 hexadecimal characters) to preserve this information.
- The hexadecimal representation of the hash is
 - H[0]H[1] → checksum
 - H[2]H[3] → L value
 - H[4] → Q1 ratio
 - H[5] → Q2 ratio
 - H[6] .. H[69] the binary representations of the 128 buckets (using the method from the previous slide) turned into hex characters

Algorithm to determine the hash

- The L value
- The input for the L value is the length of the original document (len)

```
If len <= 656
    i ←  $\log(\text{len})/\log(1.5)$ 
else
    if len <= 3199
        i ←  $\log(\text{len})/\log(1.3) - 8.72777$ 
    else
        i ←  $\log(\text{len})/\log(1.1) - 62.5472$ 
i ← i MOD 256
return i
```

Algorithm to determine TLSH

- The Q ratio values

$q1 \leftarrow$ the 32nd smallest number in bucket[0..127]

$q2 \leftarrow$ the 64th smallest number in bucket[0..127]

$q3 \leftarrow$ the 96th smallest number in bucket[0..127]

$q1_ratio \leftarrow (q1 * 100 / q3) \text{ MOD } 16$

$q2_ratio \leftarrow (q2 * 100 / q3) \text{ MOD } 16$

Calculating the distance between 2 hashes

Define a `mod_diff(X, Y, R)` function between two values, X and Y, according to a range R.

X and Y are values in the range [0, .. R-1]

Calculate the distance between X and Y in 2 ways

i.the difference between X and Y

ii.the difference between X and Y if you go up to R-1 and then back to 0

The `mod_diff` value is the minimum of (i) and (ii)

examples:

`mod_diff(3, 4, 16) = 1`

`mod_diff(3, 10, 16) = 7`

`mod_diff(3, 15, 16) = 4`

Calculating the distance between 2 hashes

```
// Input: t1 and t2
Const RANGE_LVALUE = 256
Const RANGE_QRATIO = 16

diff ← 0
ldiff ← mod_diff(t1.lvalue, t2.lvalue, RANGE_LVALUE);
If ldiff ≤ 1
    diff ← diff + ldiff
else
    diff ← diff + ldiff * 12;

q1diff ← mod_diff(t1.q1ratio, t2.q1ratio, RANGE_QRATIO);
If q1diff ≤ 1
    diff ← diff + q1diff
else
    diff ← diff + (q1diff-1) * 12;

q2diff ← mod_diff(t1.q2ratio, t2.q2ratio, RANGE_QRATIO);
If q2diff ≤ 1
    diff ← diff + q2diff
else
    diff ← diff + (q2diff-1) * 12;
```

Calculating the distance between 2 hashes (cont.)

```
If t1.checksum <> t2.checksum
    diff ← diff + 1

for i ← 1 to 64 {
    decode t1.H[i+5] in 4 binary values b10 b11 b12 b13
    decode t2.H[i+5] in 4 binary values b20 b21 b22 b23
    if (b10,b11) != b(20,21) {
        if (b10,b11) == (1,1) AND (b20,b21) == (0,0) diff = diff + 6
        else if (b10,b11) == (0,0) AND (b20,b21) == (1,1) diff = diff + 6
        else if (b10,b11) == (1,1) AND (b20,b21) == (0,1) diff = diff + 2
        else if (b10,b11) == (0,0) AND (b20,b21) == (1,0) diff = diff + 2
        else diff = diff + 1
    }
    // do the identical process for (b12,b13) and (b22,b23)
    ...
}

return diff
```

References

The Pearson Hash

[1] "Fast Hashing of Variable-Length Text Strings" by Peter K. Pearson
Communications of the ACM, Volume 33 Issue 6, June 1990 Pages 677-680.
<http://cs.mwsu.edu/~griffin/courses/2133/downloads/Spring11/p677-pearson.pdf>

SdHash

[2] "Data fingerprinting with similarity digests"
Vassil Roussev
Sixth IFIP WG 11.9 International Conference on Digital Forensics, Hong Kong, China, January 4-6, 2010
<http://roussev.net/pdf/2010-IFIP--sdhash-design.pdf>

Nilsimsa

[3] Source code for Nilsimsa <http://ixazon.dynip.com/~cmeclax/nilsimsa.html>

[4] "An open digest-based technique for spam detection"
E. Damiani¹, S. De Capitani di Vimercati¹, S. Paraboschi², P. Samarati
Proceedings of the 2004 international workshop on security in parallel and distributed systems. 2004.
<http://spdp.di.unimi.it/papers/pdcs04.pdf>

References (cont.)

SSDEEP

[5] "Identifying almost identical files using context triggered piecewise hashing"

Jesse Kornblum

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<http://dfrws.org/2006/proceedings/12-Kornblum.pdf>

[6] Source code for SSDEEP: <http://ssdeep.sourceforge.net/>

Comparison Paper

[7] "An evaluation of forensic similarity hashes"

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